

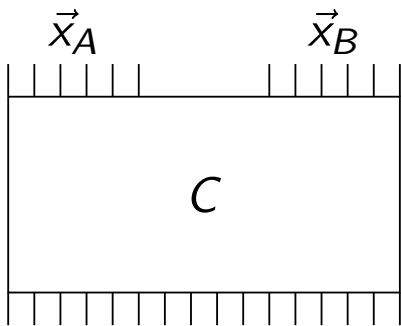
Breaking the Circuit-Size Barrier under Quasi-Polynomial LPN

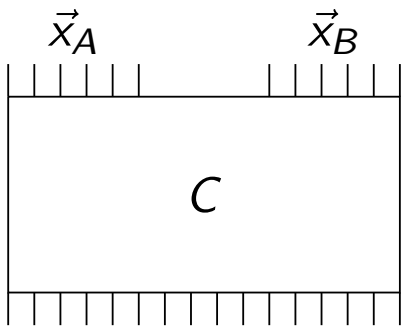
Sublinear 2-PC from LPN, the HSS Route

Geoffroy Couteau Pierre Meyer

Eurocrypt 2021







Communication

$$o(|C|)$$

Computation

$$|C|^{O(1)}$$

Correlated Randomness

FHE-based

HSS-based

Correlated Randomness

FHE-based

LWE

HSS-based

Correlated Randomness

The diagram is a large rectangle divided into three colored regions. The top region is a yellow triangle with the text 'Correlated Randomness'. The bottom-left region is a light blue trapezoid with the text 'FHE-based' and a cloud-shaped bubble containing 'LWE'. The bottom-right region is a light green trapezoid with the text 'HSS-based' and two cloud-shaped bubbles containing 'DCR' and 'LWE'. The text 'DDH' is located in the yellow triangle area.

DDH

FHE-based

LWE

DCR

HSS-based

LWE

Correlated Randomness

FHE-based

LWE

DDH

DCR

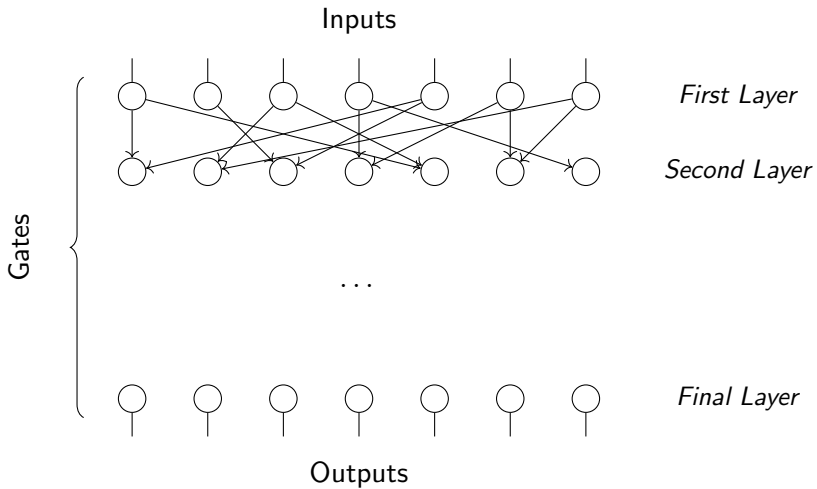
HSS-based

NEW!

LPN

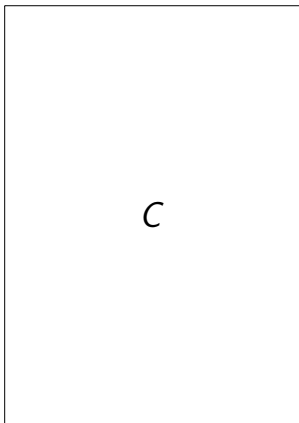
LWE

HSS-Based Sublinear 2-PC



Inputs

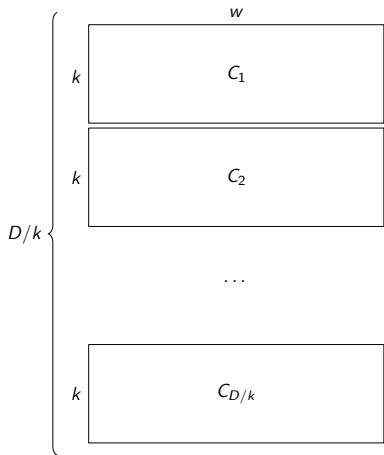
w

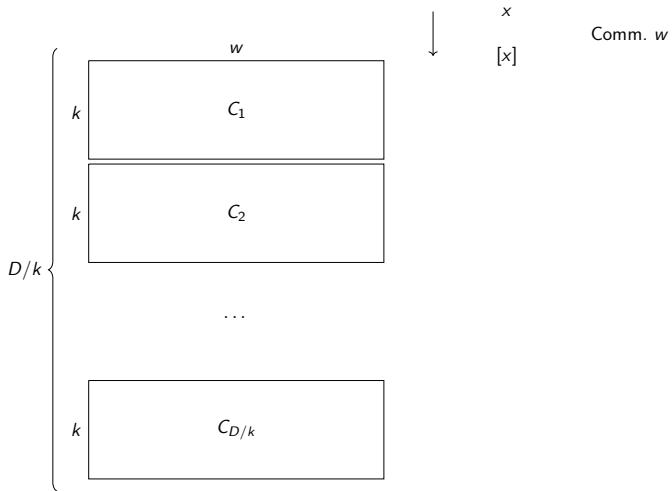


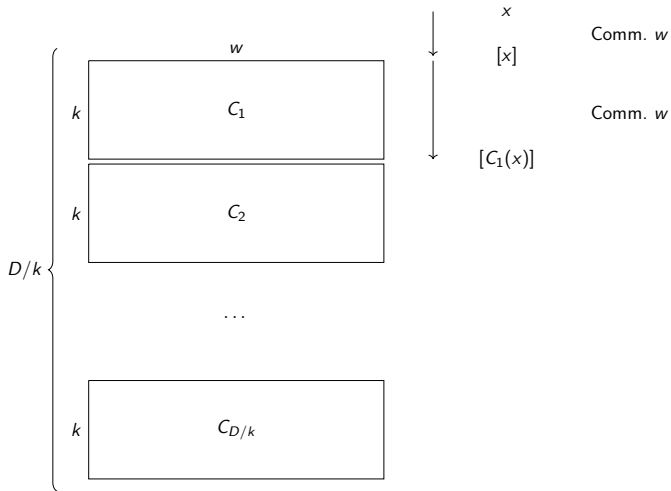
D

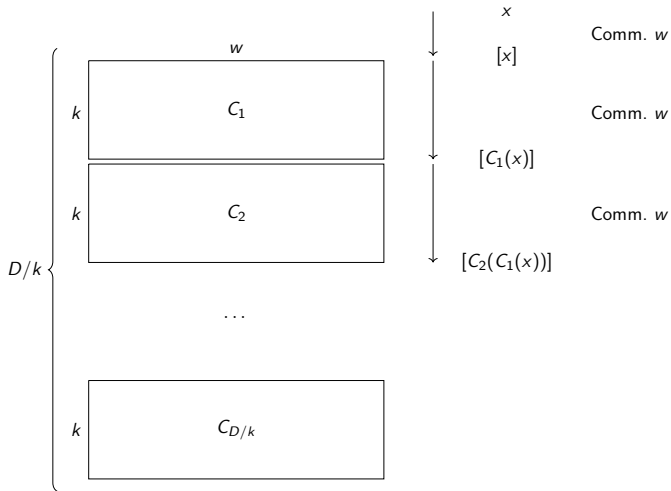
C

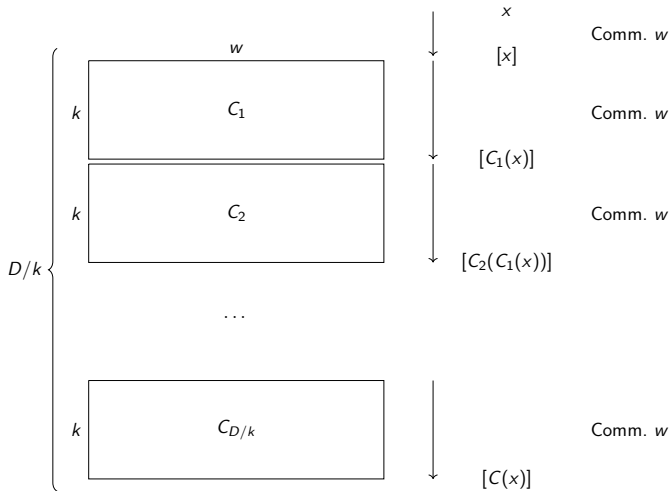
Outputs

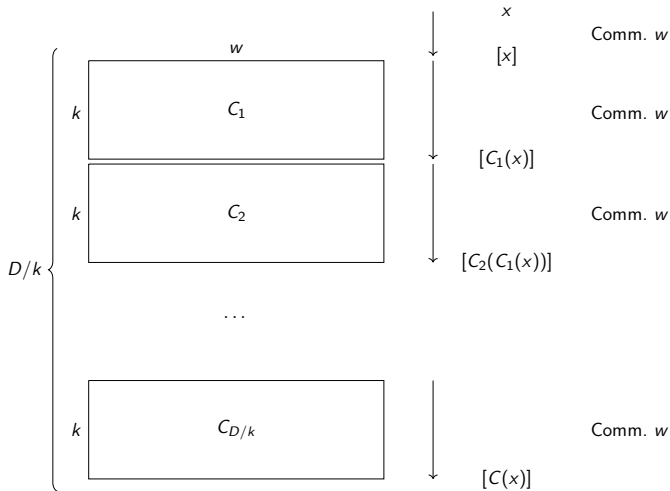












$$\text{Total: } \frac{D}{k} \cdot w = |C|/k$$

**What is the flavour of
LPN assumption?**

Quasi-Poly Learning Parity with Noise (LPN)

(Operations are over \mathbb{F}_q)

$$\text{nb of samples} \left(\begin{array}{c} \text{dimension} \\ \left(\begin{array}{c} \left[\begin{array}{c} A \\ \vdots \\ \vdots \end{array} \right] , \left[\begin{array}{c} A \\ \vdots \\ \vdots \end{array} \right] \begin{array}{c} \text{secret} \\ \downarrow \\ \left[\begin{array}{c} s \\ \vdots \\ \vdots \end{array} \right] \\ \text{error} \\ \downarrow \\ \left[\begin{array}{c} e \\ \vdots \\ \vdots \end{array} \right] \end{array} \right) + \left(\begin{array}{c} \left[\begin{array}{c} A \\ \vdots \\ \vdots \end{array} \right] , \left[\begin{array}{c} \$ \\ \vdots \\ \vdots \end{array} \right] \end{array} \right) \end{array} \right) \approx_c \left(\begin{array}{c} \left[\begin{array}{c} A \\ \vdots \\ \vdots \end{array} \right] , \left[\begin{array}{c} \$ \\ \vdots \\ \vdots \end{array} \right] \end{array} \right)$$

Quasi-Poly Learning Parity with Noise (LPN)

(Operations are over \mathbb{F}_q)

$$\text{nb of samples} \left(\left(\begin{array}{c} \text{dimension} \\ \left[\begin{array}{c} A \\ \vdots \\ \vdots \\ A \end{array} \right] \end{array} \right), \left[\begin{array}{c} A \\ \vdots \\ \vdots \\ A \end{array} \right] \begin{array}{c} \text{secret} \\ \downarrow \\ \left[\begin{array}{c} s \\ \vdots \\ \vdots \\ s \end{array} \right] \end{array} + \begin{array}{c} \text{error} \\ \downarrow \\ \left[\begin{array}{c} e \\ \vdots \\ \vdots \\ e \end{array} \right] \end{array} \right) \approx_c \left(\left[\begin{array}{c} A \\ \vdots \\ \vdots \\ A \end{array} \right], \left[\begin{array}{c} \$ \\ \vdots \\ \vdots \\ \$ \end{array} \right] \right)$$

Noise Rate:

$$r = \frac{\text{HW}(e)}{\text{nb of samples}}$$

Quasi-Poly Learning Parity with Noise (LPN)

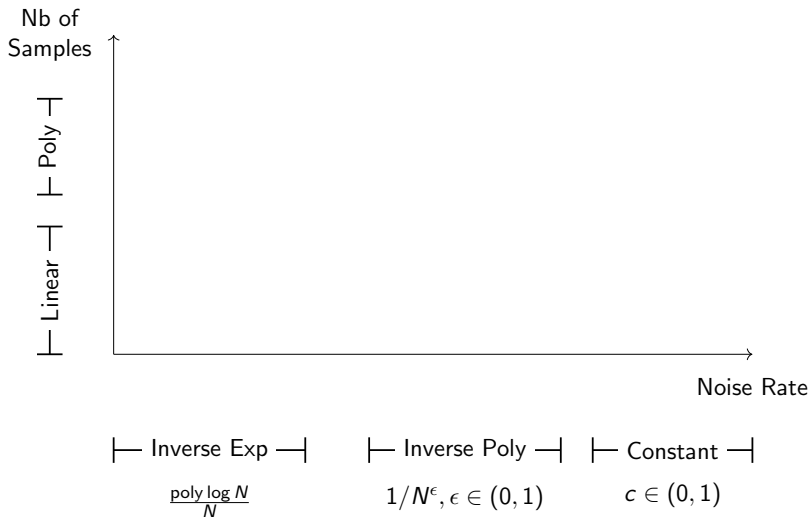
(Operations are over \mathbb{F}_q)

$$\text{nb of samples} \left(\left(\begin{array}{c} \text{dimension} \\ \left[\begin{array}{c} A \\ \left[\begin{array}{c} \text{secret} \\ s \end{array} \end{array} \right] \end{array} \right), \left[\begin{array}{c} A \\ \left[\begin{array}{c} \text{error} \\ e \end{array} \right] \end{array} \right] \right) + \left(\begin{array}{c} \left[\begin{array}{c} A \\ \left[\begin{array}{c} \$ \end{array} \right] \end{array} \right] \end{array} \right) \right) \approx_c \left(\begin{array}{c} \left[\begin{array}{c} A \\ \left[\begin{array}{c} \$ \end{array} \right] \end{array} \right] \end{array} \right)$$

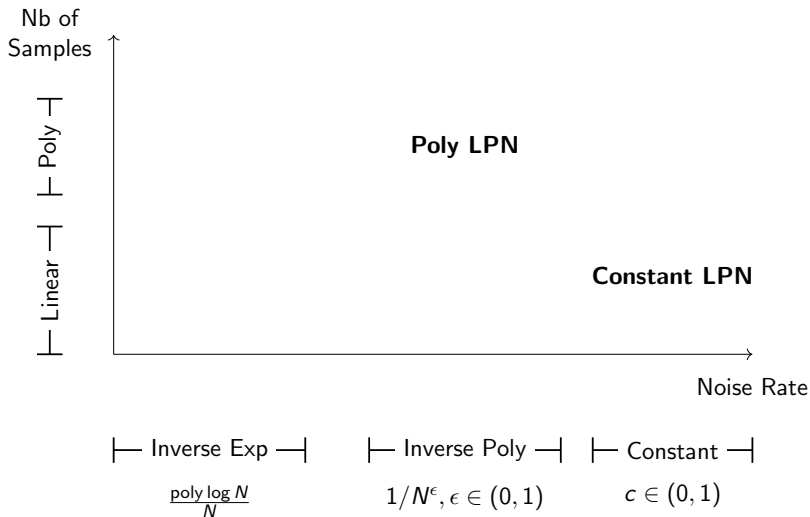
Noise Rate:

$$r = \frac{\text{HW}(e)}{\text{nb of samples}}$$

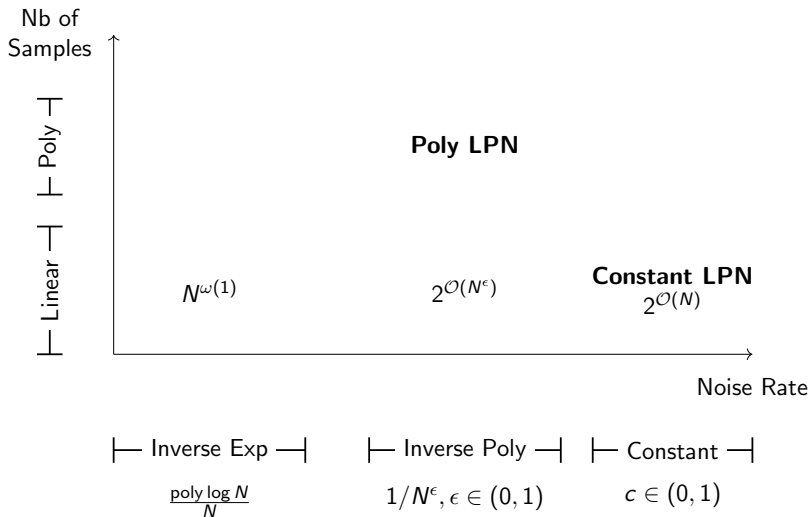
Adv. Runtime: $\lambda^{2 \log \lambda}$
Dimension: $N = \lambda^{\log \lambda}$
Nb of Samples: $2N$
Noise Rate: λ/N



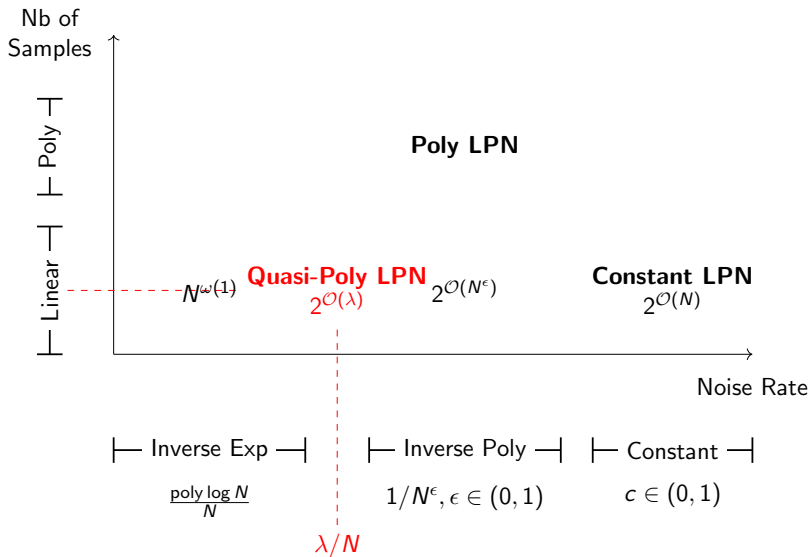
N : Dimension



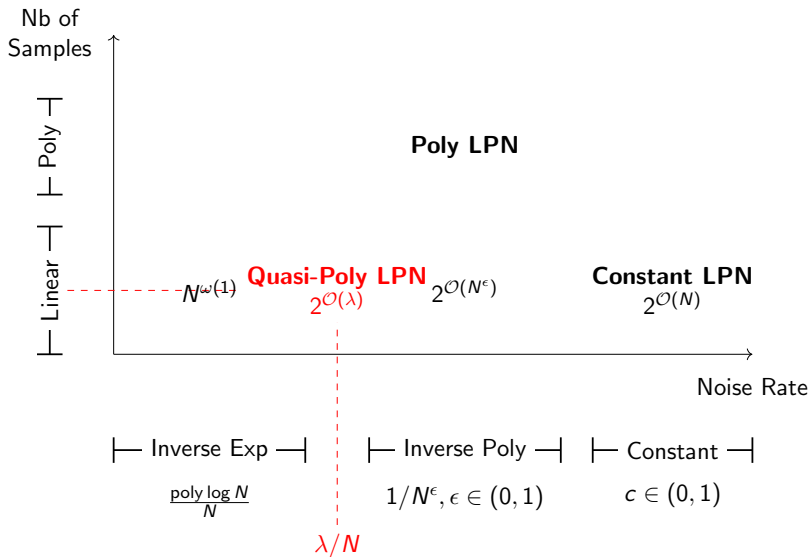
N : Dimension



N : Dimension



N : Dimension



N : Dimension

We only require security against $2^{\log^2 \lambda}$:)

A Tutorial On Naming a Paper for Optimal Metrics

Breaking the Circuit-Size Barrier for Secure Computation under Quasi-Polynomial LPN
Couteau, M. (Eurocrypt'21)

Breaking the Circuit-Size Barrier for Secure Computation under **Quasi-Polynomial LPN**
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$$\text{Normalised Citation Count} = \frac{\text{Number of Citations}}{\text{Age of Paper}}$$

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Normalized Top-100 Crypto Papers

This webpage is an attempt to assemble a ranking of top-cited papers from the area of cryptography. The ranking has been created based on citations of papers published at top cryptography conferences. More details are [available here](#).

Absolute citations are not necessarily a good indicator for the impact of a paper, as the number of citations usually grows with the age of a paper. The following list shows an alternative ranking, where the citations are normalized by the age of each paper.



93 cites??

Top 100 papers normalized by age ▾

Geoffroy Couteau and Pierre Meyer:

Breaking the Circuit Size Barrier for Secure Computation Under Quasi-Polynomial LPN.
Conference on the Theory and Applications of Cryptographic Techniques (EUROCRYPT), 2021

93 cites at [Google Scholar](#) | 3691% above average of year | Last visited: Sep-2021 | Paper: DOI

1

Cynthia Dwork, Frank McSherry, Kobbi Nissim, and Adam D. Smith:

Calibrating Noise to Sensitivity in Private Data Analysis.
Theory of Cryptography Conference (TCC), 2006

5248 cites at [Google Scholar](#) | 3144% above average of year | Last visited: Sep-2021 | Paper: DOI

2

Paul C. Kocher, Joshua Jaffe, and Benjamin Jun:

Differential Power Analysis.
International Cryptology Conference (CRYPTO), 1999

9271 cites at [Google Scholar](#) | 3040% above average of year | Last visited: Sep-2021 | Paper: DOI

3

3694%
over avg. . . ?

Thank you!

ia.cr/2021/943